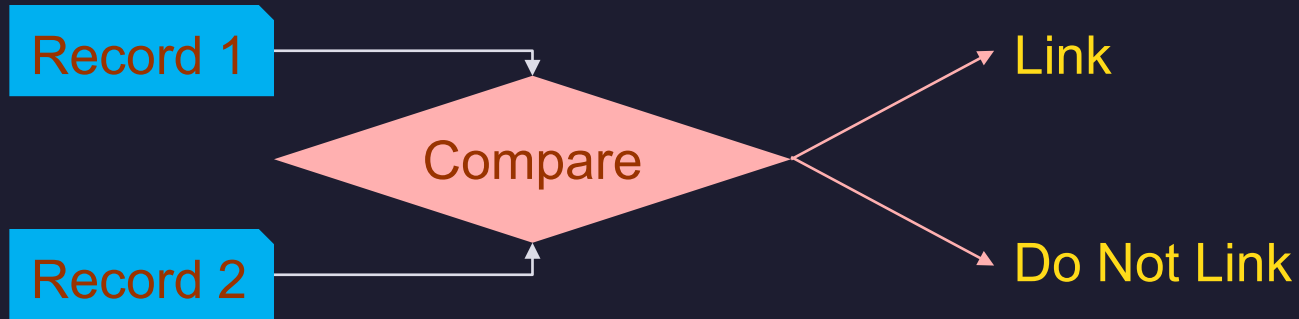

Scaling Record Linkage

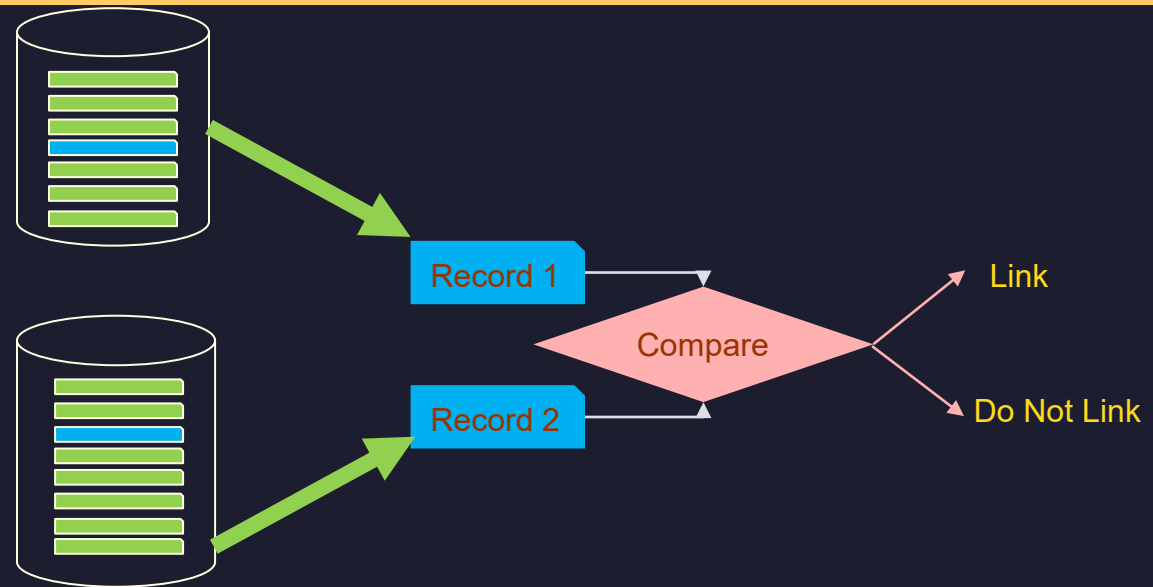
H. V. Jagadish
University of Michigan

Core Operation



- Compare operation can consider more than just the records, and can be very sophisticated (e.g. use AI methods).
- So, can be computationally expensive.

Repeated Core Operation

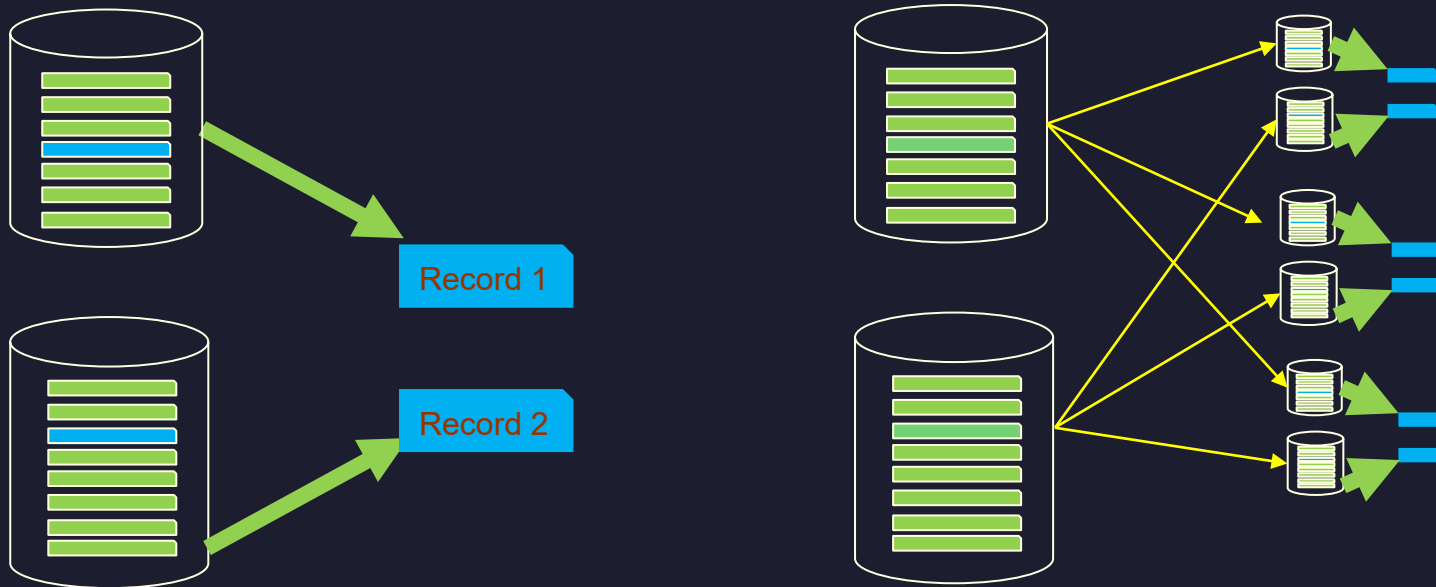


- Compare operation can be computationally expensive.
- Performed once for each pair of records.
- That is, *product* of record set sizes.
 - With 1,000 records in each record set, requires 1 million compare operations.
 - With 1 million records in each record set, requires 1 trillion compare operations.

Apply More Resources

- Limits to processing power of any one machine.
- Compare tasks can be performed in parallel.
- Be smart about distributing tasks across a bunch of processors. E.g.
 - Evenly distribute load
 - Minimize data transfer

Reduce Number of Compare

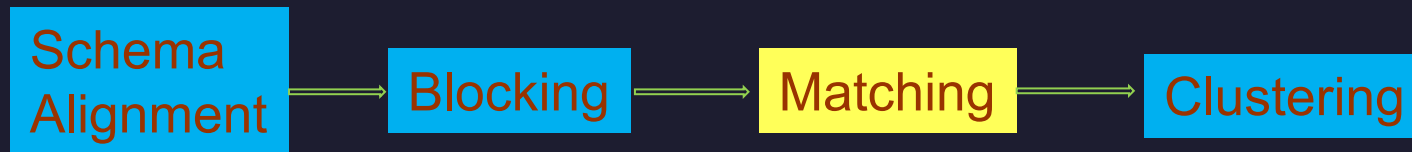


- Eliminate pairs that couldn't possibly match.
 - At least, unlikely to match.
- E.g. Use last name field as a basis to divide each record set into "blocks", and perform pairwise comparison only for records from corresponding blocks.
- With 1 million records in each record set, with blocks of 10, requires only 10 million compare operations.

Typical Workflow



Typical Workflow



- Central operation
- Typically expensive
- Typically pairwise

Typical Workflow



- Multi-source linkage
- E.g. on the web

Typical Workflow



- Divide record set into blocks.
- Must be performed cheaply.
- Based only on an individual record
 - Without comparing with others
- E.g. Use a hash to partition.
- E.g. Last name + Zipcode

Data often has errors

- E.g. Misspelt last name
- E.g. Typographical errors
- Simple blocking can put related records in different blocks, and this is not recoverable.
- Fix by having larger (potentially overlapping blocks).
 - E.g. consider letter n-grams

Limits to Basic Blocking

- Need to identify “must have” conditions.
- E.g. Changed Last Name cannot be handled.

Typical Workflow

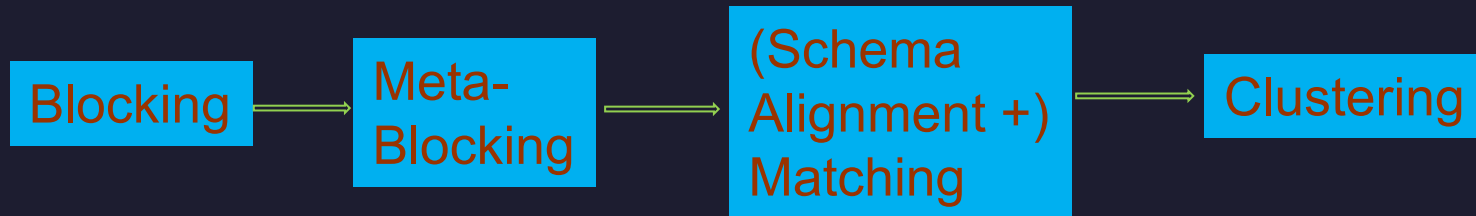


- Sources typically independent
 - Except for duplicate merging
- Corresponding attributes may have different names.
 - May even be differently structured
 - Name vs. firstname, lastname
 - Or differently expressed
 - Date formats
 - Units of measurement

Heterogeneity

- Schema alignment is hard
- Often imperfect
- Would rather address (some of it) at match time
- There may even be no schema for some records, e.g. in NL text.

Modified Workflow



Meta-blocking

- Create lots of (overlapping) blocks.
- Be generous in block creation and record assignment.
- Each record assigned to multiple blocks.
- Use meta-blocking to clean this up.

Meta-Blocking Example

p1

Name: John Abram Jr
profession: car seller
year: 1985
Addr.: Main street

p3

name1: Jon Jr
name2: Abram
birth year: 85
job: car retail
Loc: Main st.

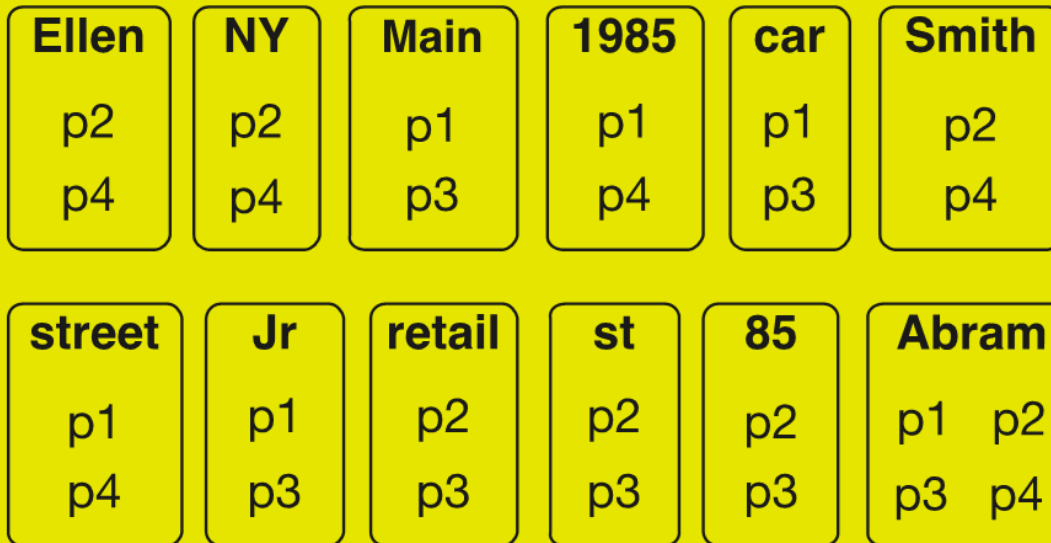
p2

FirstName: Ellen
SecondName: Smith
year: 85
occupation: retail
mail: Abram st. 30 NY

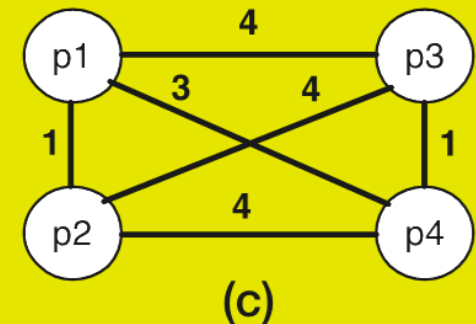
p4

full name: Ellen Smith
b. date: May 10 1985
work info: retailer
loc: Abram street NY

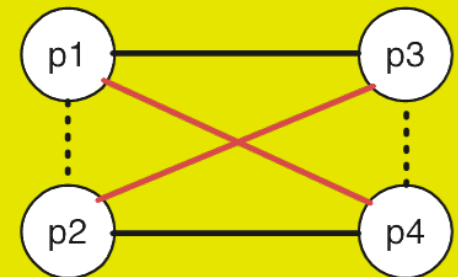
(a)



(b)

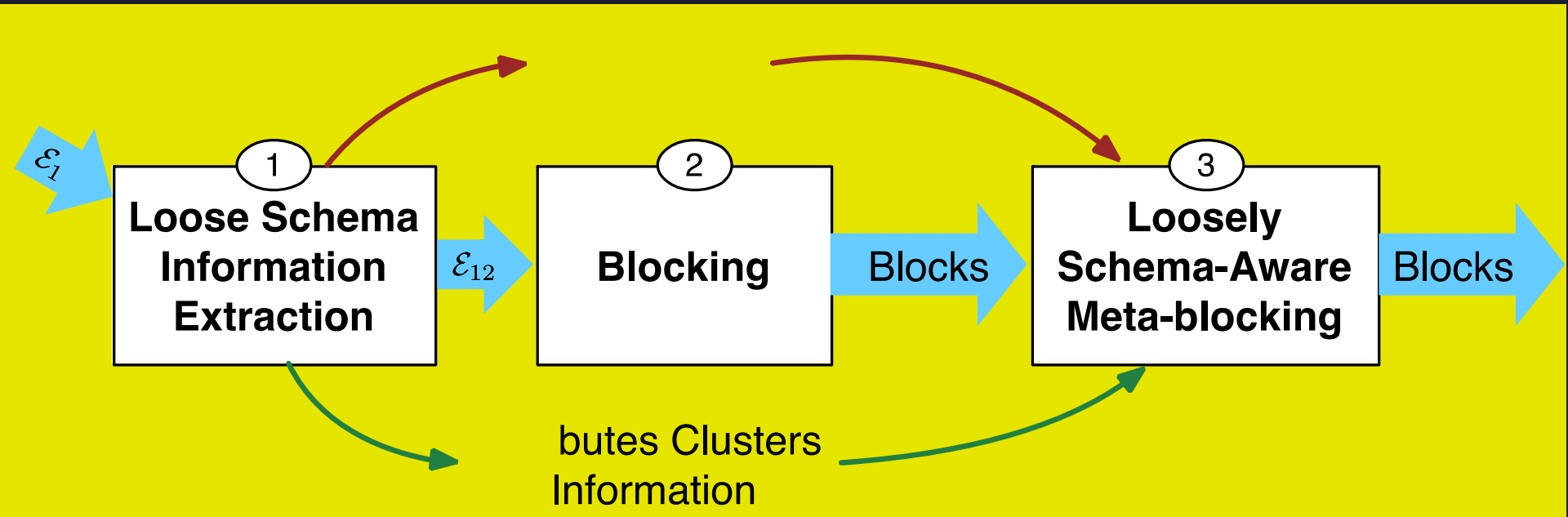


(c)

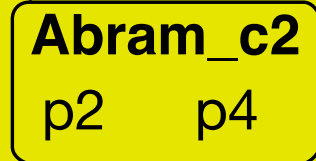


(d)

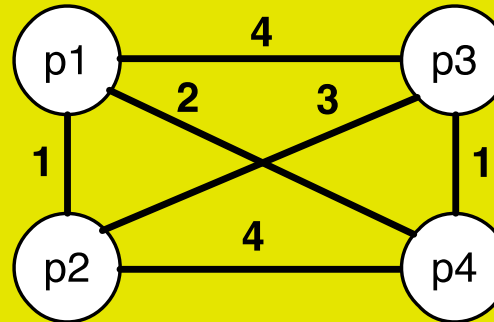
BLAST



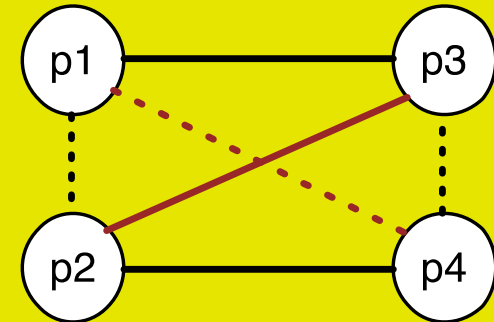
Attribute Clustering



(a)



(b)



(c)

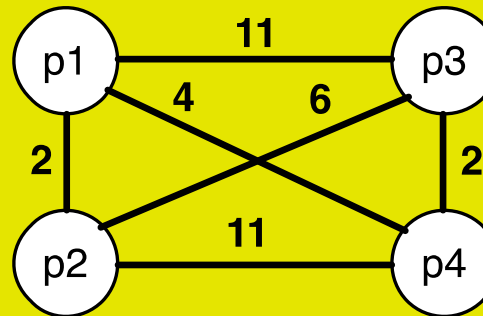
- Cluster attributes very roughly into groups.
 - Much easier than full schema alignment
- Block only for shared token in same group.

Attribute (Cluster) Entropy

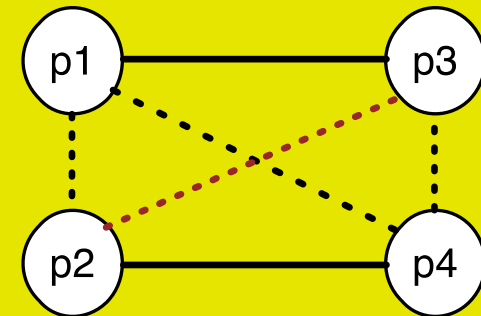
[Loose Schema Info]

Entropy cluster1 (name) = **3.5**
Entropy cluster2 (other atr.) = 2.0

(a)



(b)



(c)

- Not all attributes are equally informative
- Compute attribute entropies
 - At attribute cluster level
- Weight edges by entropy for meta-blocking

Conclusion

- Record linkage is messy.
- Many clever methods to match (not discussed today).
- But can quickly get expensive.
- Use blocking, and meta-blocking to scale.
- Also use parallelism
 - We have work to parallelize meta-blocking